

Example

matrix dimension	A_1	A_2	A_3	A_4	A_5	A_6
	4×2	2×3	3×1	1×2	2×2	2×3

Table m

M	A_1	A_2	A_3	A_4	A_5	A_6	2	3	4	5	6	
A_1	0	24	14	22	26	36	1	1	3	3	3	1
A_2		0	6	10	14	22		2	3	3	3	2
A_3			0	6	10	19			3	3	3	3
A_4				0	4	10				4	5	4
A_5					0	12					5	5
A_6						0						

Computing the Optimal Costs :-

$L=2$

$m[1,2], m[2,3], m[3,4], m[4,5], m[5,6]$

$L=3$

$m[1,3], m[2,4], m[3,5], m[4,6]$

$L=4$

$m[1,4], m[2,5], m[3,6]$

$L=5$

$m[1,5], m[2,6]$

$L=6$

$m[1,6]$

Compute each one
and store the min
value in the table m

$$m[1,2] = 4 \times 2 \times 3 = 24$$

$$m[2,3] = 2 \times 3 \times 1 = 6$$

$$m[3,4] = 3 \times 1 \times 2 = 6$$

$$m[4,5] = 1 \times 2 \times 2 = 4$$

$$m[5,6] = 2 \times 2 \times 3 = 12$$

$$m[1,3] = \min \begin{cases} m[1,1] + m[2,3] + P_0 P_1 P_3 = 0 + 6 + 4 \times 2 \times 1 = 14 \quad k=1 \\ m[1,2] + m[3,3] + P_0 P_2 P_3 = 24 + 0 + 4 \times 3 \times 1 = 36 \end{cases}$$

$$m[2,4] = \min \begin{cases} m[2,2] + m[3,4] + P_1 P_2 P_4 = 0 + 6 + 2 \times 3 \times 2 = 18 \\ m[2,3] + m[4,4] + P_1 P_3 P_4 = 6 + 0 + 4 = 10 \quad k=3 \end{cases}$$

$$m[3,5] = \min \begin{cases} m[3,3] + m[4,5] + P_2 P_3 P_5 = 0 + 4 + 6 = 10 \quad k=3 \\ m[3,4] + m[5,5] + P_2 P_4 P_5 = 6 + 0 + 12 = 18 \end{cases}$$

$$m[4,6] = \min \begin{cases} m[4,4] + m[5,6] + P_3 P_4 P_6 = 0 + 12 + 6 = 18 \\ m[4,5] + m[6,6] + P_3 P_5 P_6 = 4 + 0 + 6 = 10 \quad k=5 \end{cases}$$

$$m[1,4] = \min \begin{cases} \min[1,3] + m[2,4] + P_0 P_1 P_4 = 0 + 10 + 16 = 26 \\ \min[1,2] + m[3,4] + P_0 P_2 P_4 = 24 + 6 + 24 = 54 \\ \min[1,3] + m[4,4] + P_0 P_3 P_4 = 14 + 0 + 8 = 22 \quad k=3 \end{cases}$$

$$m[2,5] = \min \begin{cases} m[2,2] + m[3,5] + P_1 P_2 P_5 = 0 + 10 + 12 = 22 \\ m[2,3] + m[4,5] + P_1 P_3 P_5 = 6 + 4 + 4 = 14 \quad k=3 \\ m[2,4] + m[5,5] + P_1 P_4 P_5 = 10 + 0 + 8 = 18 \end{cases}$$

$$m[3,6] = \min \begin{cases} m[3,3] + m[4,6] + P_2 P_3 P_6 = 0 + 10 + 9 = 19 \quad k=3 \\ m[3,4] + m[5,6] + P_2 P_4 P_6 = 6 + 12 + 18 = 36 \\ m[3,5] + m[6,6] + P_2 P_5 P_6 = 10 + 0 + 18 = 28 \end{cases}$$

$$m[1,5] = \min \begin{cases} m[1,1] + m[2,5] + P_0 P_1 P_5 = 0 + 14 + 16 = 30 \\ m[1,2] + m[3,5] + P_0 P_2 P_5 = 24 + 10 + 24 = 58 \\ m[1,3] + m[4,5] + P_0 P_3 P_5 = 14 + 4 + 8 = 26 \quad k=3 \\ m[1,4] + m[5,5] + P_0 P_4 P_5 = 22 + 0 + 16 = 38 \end{cases}$$

$$m[2,6] = \min \begin{cases} m[2,2] + m[3,6] + P_1 P_2 P_6 = 0 + 19 + 18 = 37 \\ m[2,3] + m[4,6] + P_1 P_3 P_6 = 6 + 10 + 6 = 22 \quad k=3 \\ m[2,4] + m[5,6] + P_1 P_4 P_6 = 10 + 12 + 12 = 34 \\ m[2,5] + m[6,6] + P_1 P_5 P_6 = 14 + 0 + 12 = 26 \end{cases}$$

* * *

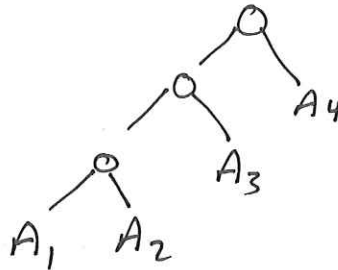
$$m[1,6] = \min \begin{cases} m[1,1] + m[2,6] + P_0 P_1 P_6 = 0 + 22 + 24 = 46 \\ m[1,2] + m[3,6] + P_0 P_2 P_6 = 24 + 19 + 36 = 79 \\ m[1,3] + m[4,6] + P_0 P_3 P_6 = 14 + 10 + 12 = 36 \quad k=3 \\ m[1,4] + m[5,6] + P_0 P_4 P_6 = 22 + 12 + 24 = 58 \\ m[1,5] + m[6,6] + P_0 P_5 P_6 = 26 + 0 + 24 = 50 \end{cases}$$

Example: Matrix Chain Multiplication Problem

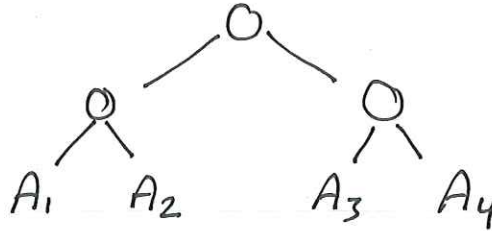
A_1 A_2 A_3 A_4
 5×4 4×6 6×2 2×7

How to multiply?

$$((A_1 \cdot A_2) \cdot A_3) \cdot A_4$$



$$(A_1 \cdot A_2) \cdot (A_3 \cdot A_4)$$



How to find the minimum number of multiplications to multiply the 4 matrices?

* Using Dynamic Programming

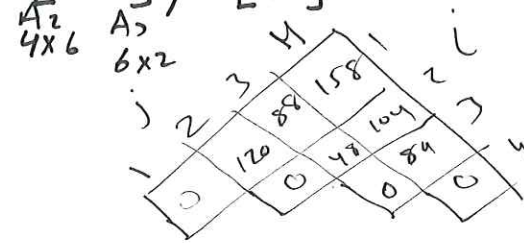
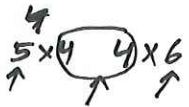
$L=1 \Rightarrow$ All zeros

✓ A_1, A_2, A_3, A_4

$L=2$

$m[1,2], m[2,3], m[3,4]$

$A_1 \cdot A_2$



m	1	2	3	4
1	0	120	88	158
2		0	48	104
3			0	84
4				0

S	1	2	3	4
1		1	1	3
2		2	2	3
3				3
4				

$L=3$

$m[1,3]$

$$(A_1 \cdot (A_2 \cdot A_3))$$

$$((A_1 \cdot A_2) \cdot A_3)$$

$$m[1,1] + m[2,3] + 5 \times 4 \times 2$$

$$m[1,2] + m[3,3] + 5 \times 6 \times 2$$

$$A_1 \cdot A_2 \cdot A_3 = 48 + 40$$

$$120 + 0 + 120$$

$$= 88$$

$$= 80$$



5×4 4×2

$5 \times 4 \times 2$



$$m[2,4] = 104 \rightarrow k=3$$

$$\begin{array}{l}
 A_2 \quad (A_3) \quad A_4 \\
 4 \times 6 \quad 6 \times 2 \quad 2 \times 7 \\
 m[2,2] + m[3,4] + 4 \times 6 \times 7 \\
 0 + 84 + 168 = 252 \\
 L=3 \\
 m[2,3] + m[4,4] + 4 \times 2 \times 7 \\
 48 + 0 + 56 \\
 104
 \end{array}$$

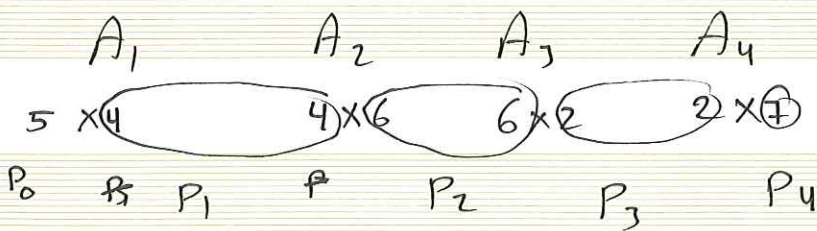
$$\begin{array}{l}
 A_1 \quad A_2 \quad A_3 \quad A_4 \\
 5 \times 4 \quad 4 \times 6 \quad 6 \times 2 \quad 2 \times 7 \\
 m[1,4] = \min \left\{ \begin{array}{l}
 m[5,1] + m[2,4] + 5 \times 4 \times 7, \quad m[1,2] + m[3,4] + \\
 0 + 104 + 140 \qquad \qquad \qquad 120 + 84 + 20 \quad (5 \times 6 \times 7)
 \end{array} \right.
 \end{array}$$

$$\begin{array}{l}
 m[1,3] + m[4,4] + 5 \times 2 \times 7 \\
 88 + 0 + 70 = 158
 \end{array}$$

k=3
s=3

Formula

$$m[i,j] = \min \left\{ m[i,k] + m[k+1,j] + P_{i-1} * P_k * P_j \right\}$$

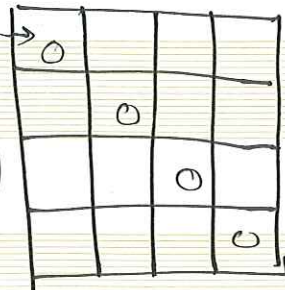


time complexity

$$\frac{n \times (n-1)}{2} = n^2 \times n \Rightarrow O(n^3)$$

we calculate only half of the table

for each value in the table, we find all possible values, (n) times

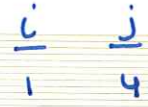


M

	0	1	2	3	4
0	0	0	0	0	0
1	0	0			
2	0		0		
3	0			0	
4	0				0

S

	0	1	2	3	4
0	0	0	0	0	0
1	0	0			
2	0		0		
3	0			0	
4	0				0



difference

$d=1$

$d=2$

$d=3$

calculate diagonally

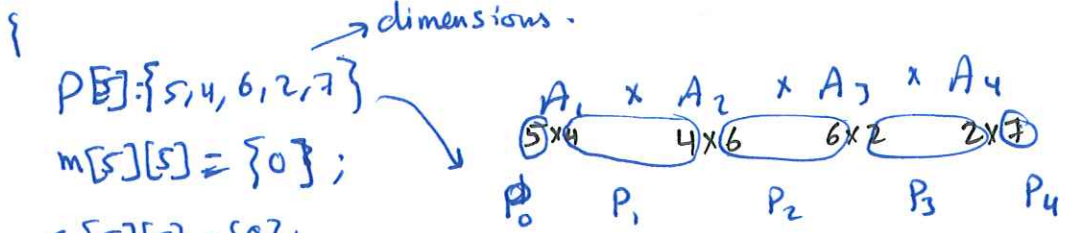
so the most outer loop becomes: for (d=1; d < n-1; d++)

$n=5$

how to find j??

$d+i$

why $d < n-1$? | the first time when the row will reach 3
 then 2 then 1. 3 is $4-1$



```
int j, min, q;
```

```
for (int d=1; d < n-1; d++)
```

```
for (i=1; i < n-d; i++)
```

```
{
```

```
    j = i+d;
```

```
    j = i+d
```

```
    min = ∞;
```

```
    for (int k=1; k <= j-1; k++)
```

```
    {
        q = m[i][k] + m[k+1][j] + P[i-1] * P[k] * P[j];
```

```
        if (q < min)
```

```
        {
            min = q;
```

```
        }
        s[i][j] = k;
```

```
    }
```

```
    m[i][j] = min;
```

```
}
```

```
}
```

```
}
```